

Optimization Algorithms

Non-Linear Mathematical Programs & KKT

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Winter 2024/25

Problem Formulation

- General **Non-linear Mathematical Program (NLP)** (constrained optimization problem):

Let $x \in \mathbb{R}^n$, $f : \mathbb{R}^n \rightarrow \mathbb{R}$, $g : \mathbb{R}^n \rightarrow \mathbb{R}^m$, $h : \mathbb{R}^n \rightarrow \mathbb{R}^l$ find

$$\min_x f(x) \quad \text{s.t.} \quad g(x) \leq 0, \quad h(x) = 0$$

- We typically assume f, g, h to be differentiable or smooth.
 - We can typically query $f, \nabla f, g, \nabla g, h, \nabla h$, optionally also $\nabla^2 f$.
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- The lecture sometimes only mentions inequality constraints g , equality constraints are analogous/easier

General approaches

- Roughly, try to somehow transform the constraint problem to
 - a series of unconstrained problems (log-barrier, AugLag, etc...)
 - a single but larger unconstrained problem (primal-dual)
 - a (series of) other constraint problems, hopefully simpler (dual, convex, SQP)

Outline

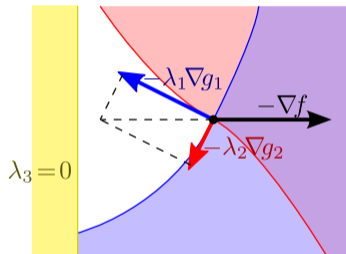
- KKT conditions of optimality
- Core methods: **log barrier**, squared penalties, **Aug. Lagrangian**
- Introduce **Lagrangian** – revisit KKT, log barrier, dual problem, primal-dual
- Further topics: Phase I, bound constraints, trust region, distributed optimization, simplex algorithm

Solving by sketching

- Sketch the following problems and identify the solution:
 - 1-dimensional: $\min x$ s.t. $\sin(x) = 0, x^2/4 - 1 \leq 0,$
 - 2-dimensional: $\min x_1$ s.t. $x^2 + y^2 - 1 \leq 0$



- *At the optimum there must be a balance between the cost gradient $-\nabla f(x)$ pulling, and the gradient of the active constraints $-\nabla g_i(x)$ pushing back*



[Our convention: “costs $f(x)$ pull, constraints $g(x), h(x)$ push back”]

Karush-Kuhn-Tucker conditions

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- **Theorem (Karush-Kuhn-Tucker conditions):** For any NLP,

$$x \text{ optimal} \quad \Rightarrow \quad \exists \lambda \in \mathbb{R}^m, \kappa \in \mathbb{R}^l \text{ s.t.}$$

$$\nabla f(x) + \sum_{i=1}^m \lambda_i \nabla g_i(x) + \sum_{j=1}^l \kappa_j \nabla h_j(x) = 0 \quad (\text{stationarity})$$

$$h(x) = 0, \quad g(x) \leq 0 \quad (\text{primal feasibility})$$

$$\lambda \geq 0 \quad (\text{dual feasibility})$$

$$\forall_i : \lambda_i g_i(x) = 0 \quad (\text{complementarity})$$

[stationarity in compact notation: $\nabla f(x) + \lambda^\top \partial_x g(x) + \kappa^\top \partial_x h(x) = 0$]

Karush-Kuhn-Tucker conditions

- **Stationarity (1st KKT):** “Force balance” of the cost pulling and the active constraints pushing back
 - Existence of dual parameters λ, κ is equivalent to

$$\nabla f(x) \in \text{span}\{\nabla g_{1..m}, \nabla h_{1..l}\}$$

- The values of λ and $\kappa \leftrightarrow$ how strongly the constraints push

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- **Complementarity (4th KKT):** “Logic of constraint activity”

- An inequality can only push at the boundary, where $g_i = 0$

- The formulation $\lambda_i g_i = 0$ very elegantly describes this logic

- The combinatorics of which constraint is active ($O(2^m)$) is a core source of difficulty of constrained optimization

- If you would apriori know which constraints are active \rightarrow inequalities become equalities \rightarrow easier